8.3. Numerical quantification

8.3.0. Overview
Claims of exemplification speak of quantity in a very limited way; but, when combined with identity, the existential quantifier can be used to express quite a variety of clearly numerical claims.

8.3.1. Else
Consider the sentence Ed signed up and someone else did, too. To analyze it as a conjunction, we need to fill out the second clause, not only by replacing did by the phrase signed up but also by making explicit an implicit reference to Ed. The full analysis would proceed as follows:

Ed signed up and someone else did, too
Ed signed up ∧ someone other than Ed signed up
Ed signed up ∧ someone other than Ed is such that (he or she signed up)
Se ∧ (3x: x is a person other than Ed) x signed up
Se ∧ (3x: x is a person and x is other than Ed) Sx
Se ∧ (3x: Px ∧ ¬ x = e) Sx
Se ∧ ∃x ((Px ∧ ¬ x = e) ∧ Sx)
P: [ _ is a person]; S: [ _ signed up]; e: Ed

That is, the function of the word else here is to restrict an existential claim by requiring that the example it claims to exist be different from a previous reference; in short, else serves to indicate a new example. The restriction of existential claims so that they claim the existence of new examples can be found not only with the word else but also, though less obviously, in a variety of quantifier phrases we have not yet attempted to analyze.

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8.3.2. Numerical quantifier phrases
So far the only numerical claims we have seen have been ones asserting or denying that a class is empty. We will now move on to a much wider group, considering claims of the three sorts

At least n Cs are such that … they …
At most n Cs are such that … they …
Exactly n Cs are such that … they …

where n may be any positive integer.

To see how to approach these quantificational claims, let us first consider existential claims regarding pairs. We looked at generalizations about pairs in 7.4.1, giving special attention to the example Not every employer and employee get along. This is the denial of generalization, so it can be understood to claim the existence of a counterexample, and we can restate it as follows to make this more explicit:

Some employer and employee do not get along.

Now we can analyze this sentence using two existential quantifiers, restricting the second by relation to the first. We would get this:

Some employer and employee do not get along

Something is such that (it and some employee of it do not get along)
3x x and some employee of x do not get along
3x some employee of x is such that (x and he or she do not get along)
3x (3y: x employs y) x and y do not get along
3x (3y: Exy) ¬ x and y get along
3x (3y: Exy) ¬ Gxy
3x 3y (Exy ∧ ¬ Gxy)
E: [_ employs _]; G: [ _ and _ get along]

The English sentence claims the existence of a pair of examples whose members are related in a certain way (as employer and employee). As with generalizations, the limitations of our notation have forced us to treat the two quantifiers asymmetrically in the symbolic form.

Now consider the sentence At least 2 things are on the agenda. It can be understood to claim the existence of a pair of examples whose members are specified to be non-identical. Following the pattern we have just used to analyze restricted existential claims concerning pairs, we can express this idea as follows:

At least 2 things are on the agenda

Something is such that (it and something else are on the agenda)
3x x and something else are on the agenda
3x something other than x is such that (x and it are on the agenda)
3x (3y: y is other than x) x and y are on the agenda
3x (3y: y is x) (x is on the agenda ∧ y is on the agenda)
3x (3y: y = x) (Nxa ∧ Nya)
3x 3y (¬ y = x ∧ A (Nxa ∧ Nya))

N: [ _ is on _ ]; a: the agenda

The quantifier phrase something else has been analyzed here before and in order to keep the vocabulary found in the quantifier phrase separate from that found in the quantified predicate of the original English sentence. But there is no need to do this, and we might have analyzed the conjunction before the second quantifier by way of an intermediate form like this:

3x x is on the agenda and so is something else

We would have ended up with the form 3x (Nxa ∧ (3y: ¬ y = x) Nya), which is equivalent to the form above by a confinement equivalence discussed in 8.1.4.

This basic idea can be extended to any quantifier phrase of the form at least n Cs. For example, at least 3 Cs can be understood to claim the existence of an example, an example different from the first, and an example different from the first two. Let us apply this idea to a case where the restrictions of non-identity are added to other specifications:

At least 3 birds are in the tree
(3x: x is a bird) x and at least 2 other birds are in the tree
(3x: Bx) (3y: y is a bird other than x) x and y and another bird are in the tree
(3x: Bx) (3y: By ∧ ¬ y = x) (3z: z is a bird other than x and y) x and y and z are in the tree
(3x: Bx) (3y: By ∧ ¬ y = x) (3z: Be ∧ (¬ z = x ∧ ¬ z = y)) (x is in the tree ∧ y is in the tree ∧ z is in the tree)
(3x: Bx) (3y: By ∧ ¬ y = x) (3z: Bz ∧ (¬ z = x ∧ ¬ z = y)) (Nxa ∧ Nyt ∧ Nzt)
B: [ _ is a bird]; N: [ _ is in _ ]; t: the tree

This can be restated in a number of different ways by using unrestricted quan-
tifiers and applying confinement principles. The following may help in thinking about the net result of the three quantifier phrases above:

\[ \exists x \forall y ((\neg x \wedge \neg y \wedge x \wedge z \equiv y) \wedge (Bx \wedge By \wedge Bz) \wedge (Nyt \wedge Nzt)) \]

That is, we assert the existence of a triple with three properties: (i) no two of its members are the same, (ii) each member is a bird, and (iii) each member is in the tree. The sentence Heinz produces at least 57 variations could be handled (in principle if not in practice) by extending the same ideas to assert the existence of a series of 57 things no two of which are the same and each of which is both a variety and produced by Heinz. If you are mathematically minded, you might try calculating the number of denied equations you would need in that case.

In the other direction, if the scopes of quantifier phrases are confined to parts of the sentence in which they bind variables, we would have instead

\[ (\exists x: Bx) (Nxt \wedge (Nyt \wedge Nzt)) \]

which might be expressed in English as Some bird is such that (i) it is in the tree and (ii) some other bird other than it is such that (a) it, too, is in the tree and (b) some bird different from both of them is in the tree also.

As a general pattern for At least \( n \) things are such that \( \ldots \) they \( \ldots \) we might use either

\[ \exists x_1 (\exists x_2 \ldots x_n ((\neg x_2 \wedge x_1) \wedge \ldots \wedge (\neg x_n \wedge x_{n-1})) \wedge \neg \forall \theta_1 \wedge \theta_2 \wedge \ldots \wedge \theta_n) \]

or

\[ \exists x_1 x_2 \ldots x_n ((\neg x_2 \wedge x_1) \wedge \ldots \wedge (\neg x_n \wedge x_{n-1})) \wedge \neg \forall \theta_1 \wedge \theta_2 \wedge \ldots \wedge \theta_n) \]

where \( \theta \) abbreviates \( \ldots \) \( \ldots \) \( \ldots \) These logical forms differ in whether the denied equations appear as restrictions on quantifiers or as conjuncts of the formula to which the quantifiers are applied. In either case, the list of denied equations should include \( \neg x_j \wedge x_i \), for each \( i > j \) where \( i, j \leq n \) —i.e., one denied equation for each pair of different variables (where requiring that the variable with the higher index appears on the left is just a systematic way of choosing one of the two ways of writing the equation). At least \( n \) Cs are such that \( \ldots \) they \( \ldots \) can be captured by adding the formulas \( x_i: C \), for each \( i \leq n \), either as restrictions on the relevant quantifiers or as further conjuncts of the quantified formula.

If we rewrite the logical forms displayed above so that quantifiers are confined to apply only to the formulas which contain variables bound to them, we would get The corresponding pattern with the quantifiers confined would be:

\[ \exists x_1 (\theta_1 x \wedge (\exists x_2 \ldots x_n ((\neg x_2 \wedge x_1) \wedge \ldots \wedge (\neg x_n \wedge x_{n-1})) \wedge \neg \forall \theta_1 \wedge \theta_2 \wedge \ldots \wedge \theta_n) \]

This says roughly, Something is such that \( \ldots \) it and so is something else \( \ldots \) and so is something else. In spite of appearances, this English sentence is not a conjunction because each use of else refers implicitly to all of the previous uses of something and cannot be treated independently from them in a separate conjunct.

We are also now in a position to analyze the other two sorts of numerical quantifier phrases mentioned earlier, for claims made using them can be restated as truth-functional compounds of claims made using at least \( n \).

At most \( n \) Cs are such that \( \ldots \) they \( \ldots \) may be paraphrased as

\( \neg \) at least \( n+1 \) Cs are such that \( \ldots \) they \( \ldots \) and

Exactly \( n \) Cs are such that \( \ldots \) they \( \ldots \) may be paraphrased as

At least \( n \) Cs are such that \( \ldots \) they \( \ldots \) at most \( n \) Cs are such that \( \ldots \) they \( \ldots \)

For example, to claim that there was at most one winner is to deny that there were at least two, and to claim that there was exactly one is to say both there was at least one and that there was at most one —i.e., it is to say that there was at least one and deny that there were at least two.

8.3.3. Exactly \( n \)

It is also possible to give a somewhat simpler symbolic representations of the quantifier phrase exactly \( n \) Cs than we get by way of truth-functional compounds of at least-\( n \) forms. Here are a couple of approaches for the case of exactly 1:

I forgot just one thing

Something is such that (I forgot it and nothing else)

\[ 3x (\neg x \wedge \neg y \wedge y = x) \wedge (x = y \wedge y \neq x) \]

I forgot just one thing

Something is such that (I forgot it and it was all I forgot)

\[ 3x (\neg x \wedge x = y) \wedge (x = \forall y (\neg y = x \wedge \neg y \neq x)) \]

\[ 3x (\neg x \wedge \neg y \wedge y = x) \wedge (x = y \wedge y \neq x) \]

I forgot just one thing

Something is such that (I forgot it and it was all I forgot)

\[ 3x (\neg x \wedge x = y) \wedge (x = \forall y (\neg y = x \wedge \neg y \neq x)) \]

And, in general, Exactly one thing is such that (\( \ldots \) it \( \ldots \)) can be analyzed as any one of the following (where \( \theta \) abbreviates \( \ldots \) \( \ldots \) \( \ldots \)):

\[ 3x (\neg y \wedge (\neg y = x) \wedge (x = \forall y (\neg y = x \wedge \neg y \neq x))) \]

\[ 3x (\neg y \wedge (\neg y = x) \wedge (x = \forall y (\neg y = x \wedge \neg y \neq x))) \]

\[ 3x \neg (\forall y (\neg y = x) \wedge \neg y \neq x) \]

The forms in rows differ only in the use of restricted and unrestricted universals, respectively, and the form columns are equivalent by the symmetry of identity and the following equivalences:

\[ (\forall y: px) x = y \equiv (\forall y: \neg \neg px) x = y \]

\[ (\forall y: \neg \neg px) x = y \equiv (\forall y: px) x = y \]

The first of these equivalences is traditionally called contraposition and that name is sometimes used for the second also. The first licenses the restatement of Only dogs barked by Everything that barked was a dog. The second would apply to the same pair of sentences when they are represented using unrestricted quantifiers and also to the restatement of The match burned only if the match burned, then oxygen was present by If the match burned, then oxygen was present.

We can also capture restricted a quantifier phrase exactly 1 C by adding restrictions to each of the two quantifiers. The following analysis of a slightly more complex example uses this sort of variation on the second pattern above:

I forgot just one number

Some number I forgot is such that (it was the numbers I forgot)

\[ (\exists x: \neg x = a \wedge x = b) \wedge (\forall x: \neg x = a \wedge \neg x = b) \]

\[ 3x (\neg x = a \wedge x = b) \wedge (\forall x: \neg x = a \wedge \neg x = b) \]

\[ 3x (\neg x = a \wedge x = b) \wedge (\forall x: \neg x = a \wedge \neg x = b) \]

And, in general, Exactly 1 C is such that (\( \ldots \) it \( \ldots \)) can be analyzed as

\[ (\exists x: \neg x = a \wedge x = b) \wedge (\forall x: \neg x = a \wedge \neg x = b) \]

The analogous variation on the first pattern would be

\[ (\exists x: x = a \wedge x = b) \wedge (\forall x: \neg x = a \wedge \neg x = b) \]

\[ (\exists x: x = a \wedge x = b) \wedge (\forall x: \neg x = a \wedge \neg x = b) \]

\[ (\exists x: x = a \wedge x = b) \wedge (\forall x: \neg x = a \wedge \neg x = b) \]

In the case of, I forgot just one number, this pattern would amount to saying Some number that I forgot is such that I forgot no other number.

The sentence There is exactly 1 C can be understood as Exactly 1 C is such that (it is) and the dummy predicate \( \_ \) can be dropped to yield the analysis

\[ (\exists x: \neg x = a \wedge \neg x = b) \wedge (\forall x: \neg x = a \wedge \neg x = b) \]

which can be understood to say Some C is such that (it is all the Cs there are).

This sort of pattern will be important for the analysis of definite descriptions in 8.4.2, but the first approach (i.e., by way of nothing else) is probably the more natural way of extending the analysis to claims of exactly \( n \) for numbers \( n > 1 \).
8.3.x. Exercise answers

1. a. If Oswald didn’t shoot Kennedy, someone else did.
   b. No one but Frank saw Sue.
   c. Ed and only Ed was awake.
   d. Everyone except Tom, Dick, and Harry arrived early.
   e. Adam and another officer thanked everyone else.
   f. At least two things went wrong.
   g. Bill spoke to at most one person.
   h. Just one thing will do.
   i. Ann saw more than one assassin.
   j. Ann saw exactly two assassins.

2. Synthesize idiomatic English sentences that express the propositions associated with the logical forms below using the intensional interpretations that follow them.

   a. Ft(ht) ∧ (Ex: x ≠ y) Ltx
   b. (Ex: Px) (3y: Py ∧ y ≠ x) Sxy
   c. (Ex: Px ∧ x ≠ m) Rex
   d. (Ex: Sx) Ox (4x: (Sx ∧ x = o) Ssx)

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e. Adam and another officer thanked everyone else
(3x: x is a officer other than Adam) Adam and x thanked everyone else
(3x: x is a officer ∧ x is other than Adam) everyone other than Adam and x is such that (Adam and x thanked him or her)
(3x: Ox ∧ ¬ x = Adam) (Wy: y is a person other than Adam and x) Adam and x both thanked y
(3x: Ox ∧ ¬ x = Adam) (Wy: y is a person ∧ y is other than Adam and x) (Adam thanked y ∧ x thanked y)
(2x: Ox ∧ ¬ x = a) (Wy: Py ∧ (¬ y = Adam ∧ ¬ y = x)) (Tat Txy)
(3x: Ox ∧ ¬ x = a) (Wy: Py ∧ (¬ y = a ∧ ¬ y = x)) (Tay Txy)
3x ((Ox ∧ ¬ x = a) ∧ Wy (Py ∧ (¬ y = a ∧ ¬ y = x))) (Tay Txy)
O: [ _ is an officer]; P: [ _ is a person]; T: [ _ thanked _]; a: Adam
or (on a different interpretation):
Adam and another officer thanked everyone else
Adam thanked everyone else ∧ an officer other than Adam thanked everyone else
everyone other than Adam is such that (Adam thanked him or her) ∧ (3x: x is a officer other than Adam) thanked everyone else
(∀ y: y is a person other than Adam) Adam thanked y ∧ (∃ x: Ox ∧ ¬ x = Adam) everyone other than x is such that (x thanked him or her)
(∀ y: Py ∧ (¬ y = x) − Adam) Tay (3x: Ox ∧ ¬ x = a) (∀ y: y is a person other than x) x thanked y
(∀ y: Py ∧ (¬ y = a) − Adam) Tay (3x: Ox ∧ ¬ x = a) (∀ y: y is a person − y)
Vy ((Py ∧ (¬ y = a) − Adam) Vy ((Py ∧ (¬ y = a) − Adam)))
The logical form produced by this second analysis is not equivalent to the one produced by the first analysis. It could be said that the first interprets else as referring to Adam and the other officer collectively while the second interprets it as referring to them individually. The latter interpretation produces a pair of generalizations each of whose domains excludes only one of Adam and the other officer rather than both together. That means that, on the second analysis, the sentence Adam and another officer thanked everyone else together with the assumption that Adam and the other officer are both people entails that they thanked each other. The second interpretation could be made more likely by stating the sentence in the form Adam and another officer each thanked everyone else.

f. At least two things went wrong
\( \exists x \: \exists y : \neg y = x \) (Tat Txy)
\( \exists x \: \exists y : \neg y = x \) (Wy: y is a person − y)
\( 3x \: 3y \: (\neg y = x) (Wx ∧ Wy) \)
W: [ _ went wrong]

g. Bill spoke to at most one person
− Bill spoke to at least two people
− at least two people are such that (Bill spoke to them)
− (∃x: x is a person) (3y: y is a person ∧ ¬ y = x) (Bill spoke to x and y)
− (∃x: Px) (3y: Py ∧ ¬ y = x) (Bill spoke to x) Bill spoke to x and y
− (∃x: Px) (3y: Py ∧ ¬ y = x) (Sbx ∧ Sby)

S: [ _ spoke to _ ]; b: Bill

h. At least one thing will do ∧ at most one thing will do
\( 3x \: x \: will \: do \: ∧ \: at \: most \: 2 \: things \: will \: do \)
\( 3x \: Dx \: x ∧ ¬ 3x (\neg y = x) (x ∧ y) \) will do
\( 3x \: Dx \: x ∧ ¬ 3x (\neg y = x) (x \: will \: do \: ∧ \: y \: will \: do) \)
\( 3x \: Dx \: x ∧ ¬ 3x (\neg y = x) (Dx ∧ Dy) \)
\( 3x \: Dx \: x ∧ ¬ 3x (\neg y = x) (Dx ∧ Dy) \)
D: [ _ will do ]
or:
\( 3x \: (x \: will \: do \: ∧ \: nothing \: other \: than \: y \: will \: do) \)
\( 3x (Dx ∧ (Wy: ¬ y = x) − y) \: will \: do \)
\( 3x (Dx ∧ (Wy: ¬ y = x) − Dy) \)
\( 3x (Dx ∧ (Wy: ¬ y = x) − Dy) \)
or:
\( 3x (x \: will \: do \: ∧ \: x \: is \: all \: that \: will \: do) \)
\( 3x (Dx ∧ (Wy: y \: will \: do) \: x = y) \)
\( 3x (Dx ∧ (Wy: y = x) − y) \)
\( 3x (Dx ∧ (Wy: − y = x) − y) \)

i. Ann saw more than one assassin
Ann saw at least two assassins
At least two assassins are such that (Ann saw them)
\( 3x \: x \: is \: an \: assassin \: ∧ \: ¬ y = x \) (Ann saw x and y)
\( 3x (Ax) (3y: Ay ∧ ¬ y = x) (Ann saw x ∧ Ann saw y) \)
\( 3x (Ax) (3y: Ay ∧ ¬ y = x) (Sax ∧ Say) \)
\( 2x (Ax ∧ 3y ((Ay ∧ ¬ y = x) ∧ (Sax ∧ Say))) \)
A: [ _ is an assassin]; S: [ _ saw _ ]; a: Ann

j. Ann saw exactly two assassins
Exactly two assassins are such that (Ann saw them)
Two assassins are such that (Ann saw them and y and no other assassins)
\( 3x (Ax) (3y: Ay ∧ ¬ y = x) (Ann saw x ∧ Ann saw y ∧ Ann saw no other assassin other than x and y) \)
\( 3x (Ax) (3y: Ay ∧ ¬ y = x) (Sax ∧ Say) ∧ no other assassin other than x and y is such that (Ann saw him or her) \)
\( 3x (Ax) (3y: Ay ∧ ¬ y = x) (Sax ∧ Say) ∧ (∀ z: z is an assassin ∧ (∼ x = z ∧ ∼ z = y)) ∧ Ann saw z) \)
\( 3x (Ax) (3y: Ay ∧ ¬ y = x) (Sax ∧ Say) ∧ (∀ z: Az (z → − x ∧ − y)) ∧ Saz) \)
\( 3x (Ax ∧ 3y ((Ay ∧ ¬ y = x) ∧ (Sax ∧ Say) ∧ ∀ z (Az (z → − x ∧ − y)) → − Saz))) \)
A: [ _ is an assassin]; S: [ _ saw _ ]; a: Ann

or:
\( 3x (Ax) (3y: Ay ∧ ¬ y = x) (Sax ∧ Say) \)
The formula (W: Ax ∧ Say) (x → y → z)) that is used here amounts to x or y are all the assassins Ann saw, for it says of any assassin that Ann saw that either x or y is that individual.

2. a. Tom found Tom’s hat ∧ (3x: x = Tom’s hat) Tom lost x
Tom found his hat ∧ (3x: x is other than Tom’s hat) Tom lost x
Tom found his hat ∧ something other than Tom’s hat is such that (Tom lost it)
Tom found his hat ∧ Tom lost something other than his hat
Tom found his hat but he lost something else

b. x: is a person (3y: y is a person ∧ ¬ y = x) x spoke to y
(3x: x is a person) (3y: y is a person ∧ y is other than x) x spoke to y
(3x: x is a person) (3y: y is a person other than x) x spoke to y
(3x: x is a person) someone other than x is such that (x spoke to him or her)
(3x: x is a person) x spoke to someone else
Someone is such that (he or she spoke to someone else)
Someone spoke to someone else

c. (3x: x is a person ∧ ¬ x = Mary) = Sam recognized x
(3x: x is a person ∧ x is other than Mary) = Sam recognized x
(3x: x is a person other than Mary) = Sam recognized x
No one other than Mary is such that (Sam recognized him or her)
Sam recognized no one other than Mary
or: Sam didn’t recognize anyone other than Mary

2. a. Tom found Tom’s hat ∧ (3x: x = Tom’s hat) Tom lost x
Tom found his hat ∧ (3x: x is other than Tom’s hat) Tom lost x
Tom found his hat ∧ something other than Tom’s hat is such that (Tom lost it)
Tom found his hat ∧ Tom lost something other than his hat
Tom found his hat but he lost something else

2. a. Tom found Tom’s hat ∧ (3x: x = Tom’s hat) Tom lost x
Tom found his hat ∧ (3x: x is other than Tom’s hat) Tom lost x
Tom found his hat ∧ something other than Tom’s hat is such that (Tom lost it)
Tom found his hat ∧ Tom lost something other than his hat
Tom found his hat but he lost something else

3x (Ax) (3y: Ay ∧ ¬ y = x) (Ann saw x ∧ Ann saw y)
3x (Ax) (3y: Ay ∧ ¬ y = x) (Sax ∧ Say)
A: [ _ is an assassin]; S: [ _ saw _ ]; a: Ann

2. a. Tom found Tom’s hat ∧ (3x: x = Tom’s hat) Tom lost x
Tom found his hat ∧ (3x: x is other than Tom’s hat) Tom lost x
Tom found his hat ∧ something other than Tom’s hat is such that (Tom lost it)
Tom found his hat ∧ Tom lost something other than his hat
Tom found his hat but he lost something else

3x (Ax) (3y: Ay ∧ ¬ y = x) (Ann saw x ∧ Ann saw y)
3x (Ax) (3y: Ay ∧ ¬ y = x) (Sax ∧ Say)
A: [ _ is an assassin]; S: [ _ saw _ ]; a: Ann

2. a. Tom found Tom’s hat ∧ (3x: x = Tom’s hat) Tom lost x
Tom found his hat ∧ (3x: x is other than Tom’s hat) Tom lost x
Tom found his hat ∧ something other than Tom’s hat is such that (Tom lost it)
Tom found his hat ∧ Tom lost something other than his hat
Tom found his hat but he lost something else

2. a. Tom found Tom’s hat ∧ (3x: x = Tom’s hat) Tom lost x
Tom found his hat ∧ (3x: x is other than Tom’s hat) Tom lost x
Tom found his hat ∧ something other than Tom’s hat is such that (Tom lost it)
Tom found his hat ∧ Tom lost something other than his hat
Tom found his hat but he lost something else

2. a. Tom found Tom’s hat ∧ (3x: x = Tom’s hat) Tom lost x
Tom found his hat ∧ (3x: x is other than Tom’s hat) Tom lost x
Tom found his hat ∧ something other than Tom’s hat is such that (Tom lost it)
Tom found his hat ∧ Tom lost something other than his hat
Tom found his hat but he lost something else

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