

Overview

Basic system

Exploitation and planning rules		Rules for closing gaps			
sentence	as a resource	as a goal	when to close	rule	
atomic sentence	none	IP	φ	QED	
negation $\neg \varphi$	CR (if φ not atomic & goal is \perp)	RAA	φ and $\neg \varphi$	\perp	Nc
conjunction $\varphi \wedge \psi$	Ext	Cnj	any	\top	ENV
disjunction $\varphi \vee \psi$	PC	PE	\perp	any	EFQ
conditional $\varphi \rightarrow \psi$	RC (if goal is \perp)	CP	$\tau_1 \rightarrow v_1, \dots, \tau_n \rightarrow v_n$	$P\tau_1 \dots \tau_n$	EC
universal $\forall x \theta x$	UI	UG	$\tau_1 \rightarrow v_1, \dots, \tau_n \rightarrow v_n$	$P\tau_1 \dots \tau_n$	DC
existential $\exists x \theta x$	PCh	NcP	$\neg P\theta_1 \dots \theta_n$	\perp	QED=
Detachment rules (optional)					
	required resources	main auxiliary		rule	
	main auxiliary				
$\neg(\varphi \wedge \psi)$	φ or ψ			MPT	
$\varphi \vee \psi$	$\neg^\pm \varphi$ or $\neg^\pm \psi$			MTP	
$\varphi \rightarrow \psi$	$\neg^\pm \varphi$	ψ		MPP	
		$\neg^\pm \psi$		MTT	

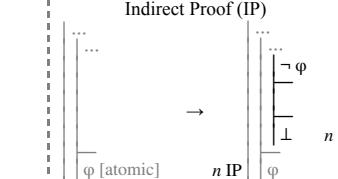
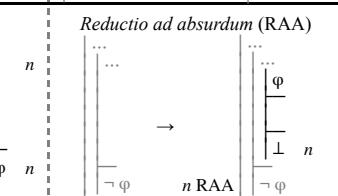
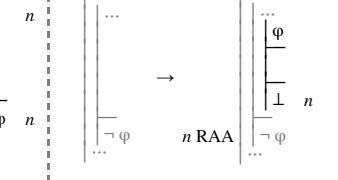
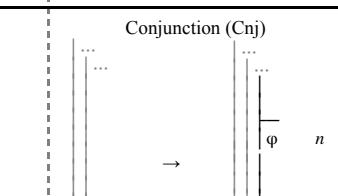
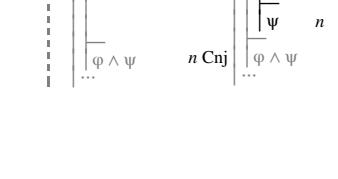
In addition, if the conditions for applying a rule are met except for differences between co-aliases, then the rule can be applied and is noted by adding “=”; QED= and Nc= are examples of this.

Additional rules

Attachment rules		Rule for lemmas	
resource to be added	rule	prerequisite	rule
$\varphi \wedge \psi$	Adj	the goal is \perp	LFR
$\neg(\varphi \wedge \psi)$			
$\varphi \vee \psi$	Wk		
$\varphi \rightarrow \psi$			
$\tau = v$	CE		
$\theta v_1 \dots v_n$	Cng		
$\exists x \theta x$	EG		

Derivation rules

Basic system

logical form	Rules for developing gaps
	as resource as goal
atomic sentence	no rule
negation $\neg \varphi$	Indirect Proof (IP) 
conjunction $\varphi \wedge \psi$	Completing the reductio (CR) 
disjunction $\varphi \vee \psi$	Reductio ad absurdum (RAA) 
conditional $\varphi \rightarrow \psi$	Modus ponendo tollens (MPT) 
universal $\forall x \theta x$	ψ [available] 
existential $\exists x \theta x$	Extraction (Ext) 
	Conjunction (Cnj) 

Rules for developing gaps		
logical form	as resource	as goal
disjunction $\varphi \vee \psi$	Proof by Cases (PC)	Proof of Exhaustion (PE)
	\vdots $\varphi \vee \psi$ \vdots \vdots \rightarrow φ χ ψ n	\vdots $\varphi \vee \psi$ n \rightarrow $\varphi \vee \psi$ n \vdots $\neg^\pm \varphi$ ψ n
	<i>Modus Tollendo Ponens (MTP)</i>	<i>OR</i>
	$\neg^\pm \varphi$ [available] \vdots $\varphi \vee \psi$ \vdots \rightarrow $\neg^\pm \varphi$ (n) \vdots $\varphi \vee \psi$ n	$\neg^\pm \varphi$ (n) \vdots $\varphi \vee \psi$ n \rightarrow $\neg^\pm \psi$ φ n
	<i>Modus Ponendo Ponens (MPP)</i>	<i>Universal Generalization (UG)</i>
	$\neg^\pm \psi$ [available] \vdots $\varphi \vee \psi$ \vdots \rightarrow $\neg^\pm \psi$ (n) \vdots $\varphi \vee \psi$ n	$\forall x \theta x$ \vdots $\rightarrow n$ UI $\theta\tau$ φ n
	<i>Modus Tollendo Tollens (MTT)</i>	<i>Universal Instantiation (UI)</i>
conditional $\varphi \rightarrow \psi$	Rejecting a Conditional (RC)	Conditional Proof (CP)
	\vdots $\varphi \rightarrow \psi$ \vdots \rightarrow φ ψ \perp n	\vdots $\varphi \rightarrow \psi$ n \rightarrow $\varphi \rightarrow \psi$ n \vdots φ ψ n
	<i>Modus Ponendo Ponens (MPP)</i>	
	φ [available] \vdots $\varphi \rightarrow \psi$ \vdots \rightarrow φ $\rightarrow \psi$ n	
existential $\exists x \theta x$	Proof by Choice (PCh)	Non-constructive Proof (NcP)
	$\exists x \theta x$ \vdots \rightarrow $\exists x \theta x$ n \vdots θa φ n	\vdots $\exists x \theta x$ n \rightarrow $\forall x \neg^\pm \theta x$ \perp $\exists x \theta x$ n

The parameter a used in UG and PCh should be new to the derivation; that is, it should appear only to the right of the scope line it labels

Rules for closing gaps (truth-functional logic)		
when to close	rule	
resources	goal	
φ	φ	Quod Erat Demonstrandum (QED)
		$\cdots \varphi [\text{available}] \cdots \rightarrow \varphi (n)$
		$\varphi \quad \varphi$
		$\varphi \quad \varphi$
		$n \text{ QED} \quad \bullet$
$\varphi \text{ and } \neg \varphi$	\perp	Non-contradiction (Nc)
		$\cdots \neg \varphi [\text{available}] \cdots \neg \varphi (n)$
		$\varphi [\text{available}] \quad \varphi (n)$
		$\perp \quad \perp$
		$n \text{ Nc} \quad \bullet$
any	\top	Ex Nihilo Verum (ENV)
		$\cdots \top \cdots \top (n)$
		$\top \quad \top$
		$n \text{ ENV} \quad \bullet$
\perp	any	Ex Falso Quodlibet (EFQ)
		$\cdots \perp \cdots \perp (n)$
		$\perp \quad \perp$
		$n \text{ EFQ} \quad \bullet$

Rules for closing gaps (equations)		
when to close	rule	
co-aliases	resources	goal
$\tau = v$	any	Equated Co-aliases (EC)
		$\cdots [\tau \text{ and } v \text{ are co-aliases}] \cdots \tau = v (n)$
		$\tau = v \quad \tau = v$
		$n \text{ EC} \quad \bullet$
$\tau = v$	$\neg \tau = v$	Distinguished Co-aliases (DC)
		$\cdots [\tau \text{ and } v \text{ are co-aliases}] \cdots \neg \tau = v (n)$
		$\neg \tau = v \quad \neg \tau = v$
		$n \text{ DC} \quad \bullet$
$\tau_1 = v_1, \dots, \tau_n = v_n$	$P\tau_1 \dots \tau_n$	QED given equations (QED=)
		$\cdots [\text{have co-alias relations:}] \cdots [\text{have co-alias relations:}]$
		$\tau_1 = v_1, \dots, \tau_n = v_n \quad \tau_1 = v_1, \dots, \tau_n = v_n$
		$P\tau_1 \dots \tau_n \quad (n)$
		$n \text{ QED=} \quad \bullet$
$\tau_1 = v_1, \dots, \tau_n = v_n$	$\neg Pv_1 \dots v_n$	Non-contradiction given equations (Nc=)
		$\cdots [\text{have co-alias relations:}] \cdots [\text{have co-alias relations:}]$
		$\tau_1 = v_1, \dots, \tau_n = v_n \quad \tau_1 = v_1, \dots, \tau_n = v_n$
		$\neg Pv_1 \dots v_n \quad (n)$
		$n \text{ Nc=} \quad \bullet$

Additional rules (not guaranteed to be progressive)

Attachment rules			
what is required	added resource	rule	
φ and ψ are both available	$\varphi \wedge \psi$	Adjunction (Adj)	$\frac{\dots \varphi \text{ [available]} \dots \psi \text{ [available]} \dots}{\varphi \wedge \psi} \rightarrow n \text{ Adj} X$
$\neg^\pm \varphi$ or $\neg^\pm \psi$ is available	$\neg(\varphi \wedge \psi)$	Weakening (Wk)	$\frac{\dots \neg^\pm \varphi \text{ [available]} \dots}{\neg(\varphi \wedge \psi)} \rightarrow n \text{ Wk} X$ $\frac{\dots \neg^\pm \psi \text{ [available]} \dots}{\neg(\varphi \wedge \psi)} \rightarrow n \text{ Wk} X$
φ or ψ is available	$\varphi \vee \psi$		$\frac{\dots \varphi \text{ [available]} \dots}{\varphi \vee \psi} \rightarrow n \text{ Wk} X$ $\frac{\dots \psi \text{ [available]} \dots}{\varphi \vee \psi} \rightarrow n \text{ Wk} X$
$\neg^\pm \varphi$ or ψ is available	$\varphi \rightarrow \psi$		$\frac{\dots \neg^\pm \varphi \text{ [available]} \dots}{\varphi \rightarrow \psi} \rightarrow n \text{ Wk} X$ $\frac{\dots \psi \text{ [available]} \dots}{\varphi \rightarrow \psi} \rightarrow n \text{ Wk} X$

Attachment rules				
what is required	added resource	rule		
τ and v are co-aliases	$\tau = v$	Co-alias Equation (CE) \dots $[\tau \text{ and } v \text{ are co-aliases}]$ \dots \dots $\vdash \varphi$ \dots	$\rightarrow n \text{ CE}$ \dots $[\tau \text{ and } v \text{ are co-aliases}]$ \dots \dots $\vdash \varphi$ \dots	
have co-alias relations $\tau_1 = v_1, \dots, \tau_n = v_n$ and $\theta\tau_1 \dots \tau_n$ is available	$\theta v_1 \dots v_n$	Congruence (Cng) \dots $[\text{have co-alias relations:}]$ $\tau_1 = v_1, \dots, \tau_n = v_n$ \dots $\theta\tau_1 \dots \tau_n$ \dots $\vdash \varphi$ \dots	\dots $[\text{have co-alias relations:}]$ $\tau_1 = v_1, \dots, \tau_n = v_n$ \dots $\theta\tau_1 \dots \tau_n$ (n) \dots $\theta v_1 \dots v_n$ \dots $\vdash \varphi$ \dots	
$\theta\tau$ is available	$\exists x \, \theta x$	Existential Generalization (EG) \dots $\theta\tau$ \dots $\vdash \varphi$ \dots	\dots $\theta\tau$ (n) \dots $\exists x \, \theta x$ X \dots	
Rule for lemmas				
prerequisite	rule			
the goal is \perp	Lemma for Reductio (LFR) \dots \dots $\vdash \perp$ \rightarrow $n \text{ LFR}$		\dots \dots $\vdash \varphi$ n $\vdash \varphi$ n $\vdash \perp$ n $\vdash \perp$ \dots	