

Appendices

Appendix A. Reference

A.o. Overview

A.1. Basic concepts

Definitions of entailment and related ideas

A.2. Logical forms

Forms expressed using one or more logical constants together with symbolic and English notation or readings

A.3. Truth tables

Tables that stipulate the meaning of the constants of truth-functional logic

A.4. Derivation rules

A guide to the use of derivation rules with links to the rules themselves

[Appendix A in pdf format](#)

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A.1. Basic concepts

<i>Concept</i>	<i>Negative definition</i>	<i>Positive definition</i>
φ is <i>entailed</i> by Γ $\Gamma \Rightarrow \varphi$	There is no logically possible world in which φ is false while all members of Γ are true.	φ is true in every logically possible world in which all members of Γ are true.
φ and ψ are <i>(logically) equivalent</i> $\varphi \Leftrightarrow \psi$	There is no logically possible world in which φ and ψ have different truth values.	φ and ψ have the same truth value as each other in every logically possible world.
φ is a <i>tautology</i> $\Rightarrow \varphi$ (or $\top \Rightarrow \varphi$)	There is no logically possible world in which φ is false.	φ is true in every logically possible world.
φ is <i>inconsistent with</i> Γ $\Gamma, \varphi \Rightarrow$ (or $\Gamma, \varphi \Rightarrow \perp$)	There is no logically possible world in which φ is true while all members of Γ are true.	φ is false in every logically possible world in which all members of Γ are true.
Γ is <i>inconsistent</i> $\Gamma \Rightarrow$ (or $\Gamma \Rightarrow \perp$)	There is no logically possible world in which all members of Γ are true.	In every logically possible world, at least one member of Γ is false.
φ is <i>absurd</i> $\varphi \Rightarrow$ (or $\varphi \Rightarrow \perp$)	There is no logically possible world in which φ is true.	φ is false in every logically possible world.
Σ is <i>rendered exhaustive</i> by Γ $\Gamma \Rightarrow \Sigma$	There is no logically possible world in which all members of Σ are false while all members of Γ are true.	At least one member of Σ is true in each logically possible world in which all members of Γ are true

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A.2. Logical forms

Forms for which there is symbolic notation

	Symbolic notation	English notation or English reading	
Negation	$\neg \phi$	not ϕ	
Conjunction	$\phi \wedge \psi$	both ϕ and ψ	(ϕ and ψ)
Disjunction	$\phi \vee \psi$	either ϕ or ψ	(ϕ or ψ)
The conditional	$\phi \rightarrow \psi$	if ϕ then ψ	(ϕ implies ψ)
	$\psi \leftarrow \phi$	yes ψ if ϕ	(ψ if ϕ)
Identity	$\tau = \upsilon$	τ is υ	
Predication	$\theta \tau_1 \dots \tau_n$	θ fits τ_1, \dots, τ_n	A series of terms τ_1, \dots, τ_n can be read (series) τ_1, \dots, τ_n
Compound term	$Y \tau_1 \dots \tau_n$	Y of τ_1, \dots, τ_n Y applied to τ_1, \dots, τ_n	an' τ_n (using the contraction an' to distinguish this use of and from its use in conjunction and adding series when necessary to avoid ambiguity)
Predicate abstract	$[\phi]_{x_1 \dots x_n}$	what ϕ says of $x_1 \dots x_n$	
Functor abstract	$[\tau]_{x_1 \dots x_n}$	τ for $x_1 \dots x_n$	
Universal quantification	$\forall x \theta x$	forall x θx everything, x , is such that θx	
Restricted universal	$(\forall x: \rho x) \theta x$	forall x st ρx θx everything, x , such that ρx is such that θx	
Existential quantification	$\exists x \theta x$	forsome x θx something, x , is such that θx	
Restricted existential	$(\exists x: \rho x) \theta x$	forsome x st ρx θx something, x , such that ρx is such that θx	
Definite description	$!x \rho x$	the x st ρx the thing, x , such that ρx	

Some paraphrases of other forms

Truth-functional compounds

neither ϕ nor ψ	$\neg (\phi \vee \psi)$ $\neg \phi \wedge \neg \psi$
ψ only if ϕ	$\neg \psi \leftarrow \neg \phi$
ψ unless ϕ	$\psi \leftarrow \neg \phi$

Generalizations

All Cs are such that (... they ...)	$(\forall x: x \text{ is a } C) \dots x \dots$
No Cs are such that (... they ...)	$(\forall x: x \text{ is a } C) \neg \dots x \dots$
Only Cs are such that (... they ...)	$(\forall x: \neg x \text{ is a } C) \neg \dots x \dots$
with: among Bs	add to the restriction: $x \text{ is a } B$
except Es	$\neg x \text{ is an } E$
other than τ	$\neg x = \tau$

Numerical quantifier phrases

At least 1 C is such that (... it ...)	$(\exists x: x \text{ is a } C) \dots x \dots$
At least 2 Cs are such that (... they ...)	$(\exists x: x \text{ is a } C) (\exists y: y \text{ is a } C \wedge \neg y = x) (\dots x \dots \wedge \dots y \dots)$
Exactly 1 C is such that (... it ...)	$(\exists x: x \text{ is a } C) (\dots x \dots \wedge (\forall y: y \text{ is a } C \wedge \neg y = x) \neg \dots y \dots)$ or $(\exists x: x \text{ is a } C) (\dots x \dots \wedge (\forall y: y \text{ is a } C \wedge \dots y \dots) x = y)$

Definite descriptions (on Russell's analysis)

The C is such that (... it ...)	$(\exists x: x \text{ is a } C \wedge (\forall y: \neg y = x) \neg y \text{ is a } C) \dots x \dots$ or $(\exists x: x \text{ is a } C \wedge (\forall y: y \text{ is a } C) x = y) \dots x \dots$
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A.3. Truth tables

<i>Tautology</i>	<i>Absurdity</i>	<i>Negation</i>
$\frac{\top}{\top}$	$\frac{\perp}{\text{F}}$	$\frac{\varphi}{\top} \mid \neg \varphi$ $\frac{\text{F}}{\text{F}} \mid \text{T}$
<i>Conjunction</i>	<i>Disjunction</i>	<i>Conditional</i>
$\frac{\varphi \ \psi}{\top \ \top} \mid \varphi \wedge \psi$ $\frac{\text{T} \ \text{T}}{\text{T}} \mid \text{T}$ $\frac{\text{T} \ \text{F}}{\text{F}} \mid \text{F}$ $\frac{\text{F} \ \text{T}}{\text{F}} \mid \text{F}$ $\frac{\text{F} \ \text{F}}{\text{F}} \mid \text{F}$	$\frac{\varphi \ \psi}{\top \ \top} \mid \varphi \vee \psi$ $\frac{\text{T} \ \text{T}}{\text{T}} \mid \text{T}$ $\frac{\text{T} \ \text{F}}{\text{T}} \mid \text{T}$ $\frac{\text{F} \ \text{T}}{\text{T}} \mid \text{T}$ $\frac{\text{F} \ \text{F}}{\text{F}} \mid \text{F}$	$\frac{\varphi \ \psi}{\top \ \top} \mid \varphi \rightarrow \psi$ $\frac{\text{T} \ \text{T}}{\text{T}} \mid \text{T}$ $\frac{\text{T} \ \text{F}}{\text{F}} \mid \text{F}$ $\frac{\text{F} \ \text{T}}{\text{T}} \mid \text{T}$ $\frac{\text{F} \ \text{F}}{\text{T}} \mid \text{T}$

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A.4. Derivation rules

Basic system

Rules for developing gaps			Rules for closing gaps			
for resources		for goals	when to close			rule
atomic sentence		IP	co-aliases	resources	goal	
				φ	φ	<u>QED</u>
negation $\neg \varphi$	<u>CR</u> (if φ not atomic & goal is \perp)	<u>RAA</u>	φ and $\neg \varphi$		\perp	<u>Nc</u>
					\top	<u>ENV</u>
			\perp			<u>EFQ</u>
conjunction $\varphi \wedge \psi$	<u>Ext</u>	<u>Cnj</u>	$\tau - u$		$\tau = u$	<u>EC</u>
					\perp	<u>DC</u>
disjunction $\varphi \vee \psi$	<u>PC</u>	<u>PE</u>	$\tau - u$	$\neg \tau = u$	\perp	<u>QED=</u>
					$\tau_1 - u_1, \dots, \tau_n - u_n$	$\text{P}\tau_1 \dots \tau_n$
conditional $\varphi \rightarrow \psi$	<u>RC</u> (if goal is \perp)	<u>CP</u>			$\text{P}u_1 \dots u_n$	<u>Nc=</u>
					$\tau_1 - u_1, \dots, \tau_n - u_n$	$\neg \text{P}u_1 \dots u_n$
universal $\forall x \theta x$	<u>UI</u>	<u>UG</u>				
existential $\exists x \theta x$	<u>PCh</u>	<u>NcP</u>				

Detachment rules (optional)

required resources		rule
main	auxiliary	
		$\frac{\varphi}{\neg^{\pm} \psi} \quad \text{MPP}$
		$\frac{\varphi \rightarrow \psi}{\neg^{\pm} \psi} \quad \text{MTT}$
		$\frac{\varphi \vee \psi}{\neg^{\pm} \varphi \text{ or } \neg^{\pm} \psi} \quad \text{MTP}$
		$\frac{\neg(\varphi \wedge \psi)}{\varphi \text{ or } \psi} \quad \text{MPT}$

In addition, if the conditions for applying a rule are met except for differences between co-aliases, then the rule can be applied and is notated by adding "="; QED= and Nc= are examples of this.

Additional rules (not guaranteed to be progressive)

Attachment rules	
added resource	rule
$\varphi \wedge \psi$	<u>Adj</u>
$\varphi \rightarrow \psi$	<u>Wk</u>
$\varphi \vee \psi$	<u>Wk</u>
$\neg(\varphi \wedge \psi)$	<u>Wk</u>
$\tau = u$	<u>CE</u>
$\theta u_1 \dots u_n$	<u>Cng</u>
$\exists x \theta x$	<u>EG</u>

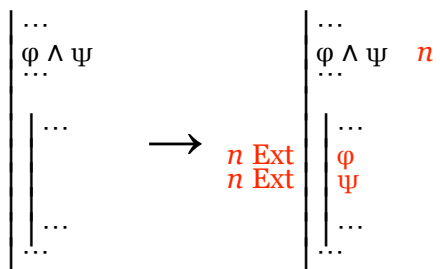
Rule for lemmas

prerequisite rule
the goal is \perp LFR

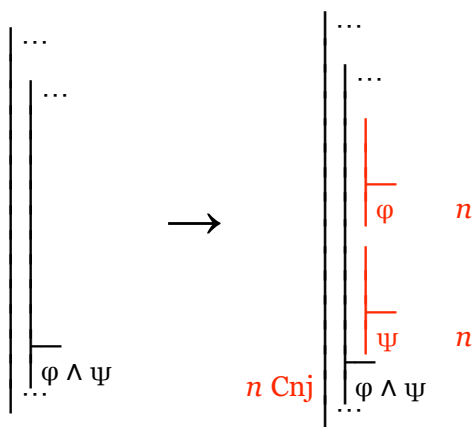
Diagrams

Rules from chapter 2

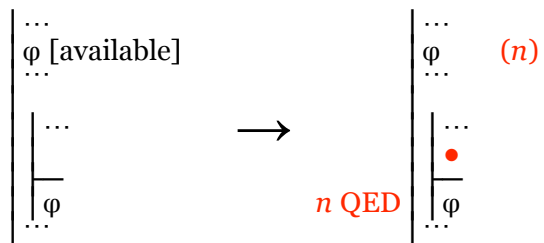
Extraction (Ext)



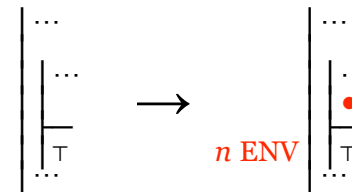
Conjunction (Cnj)



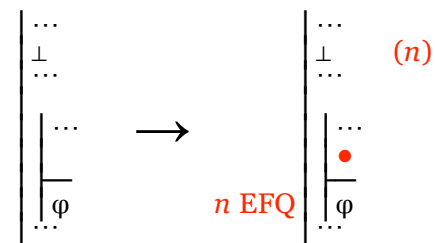
Quod Erat Demonstrandum (QED)



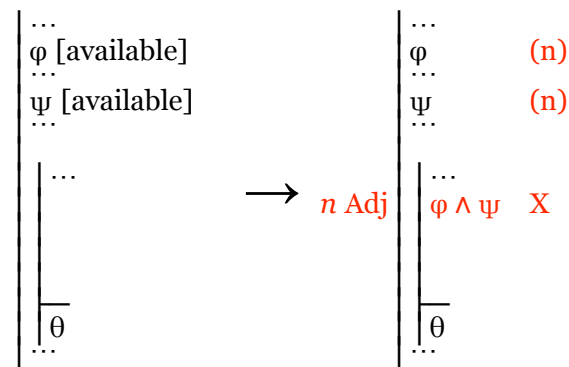
Ex Nihilo Verum (ENV)



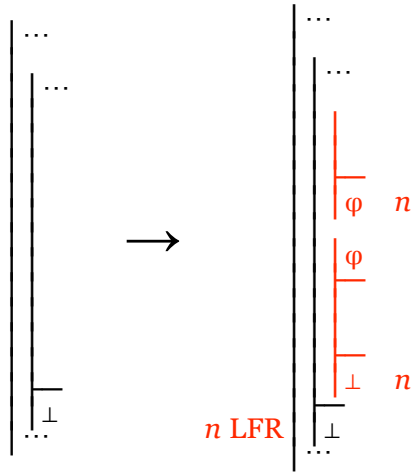
Ex Falso Quodlibet (EFQ)



Adjunction (Adj)

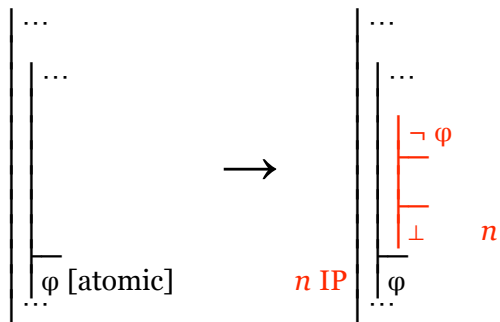


Lemma for *Reductio* (LFR)

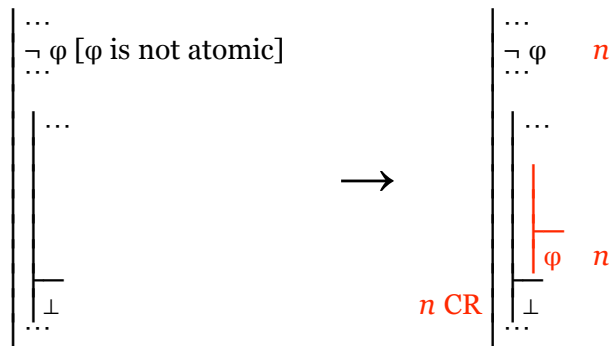


Rules from chapter 3

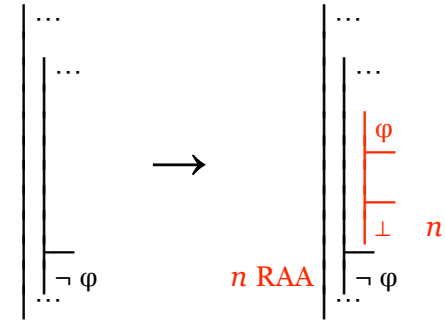
Indirect Proof (IP)



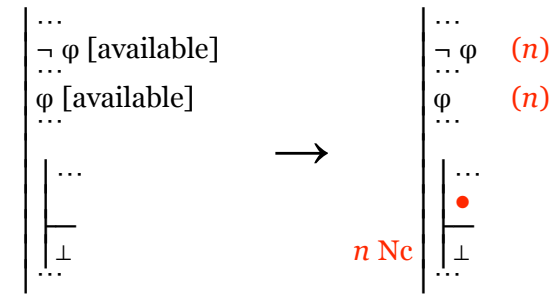
Completing the *Reductio* (CR)



Reductio ad Absurdum (RAA)

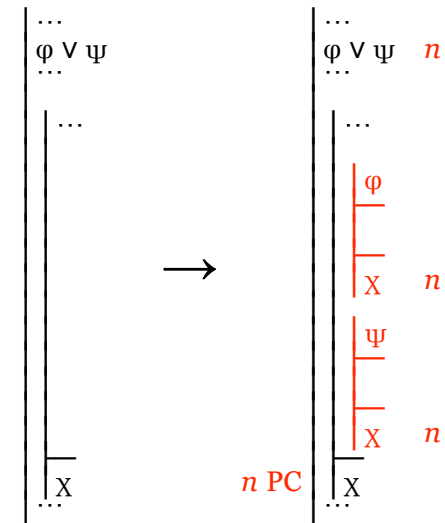


Non-contradiction (Nc)

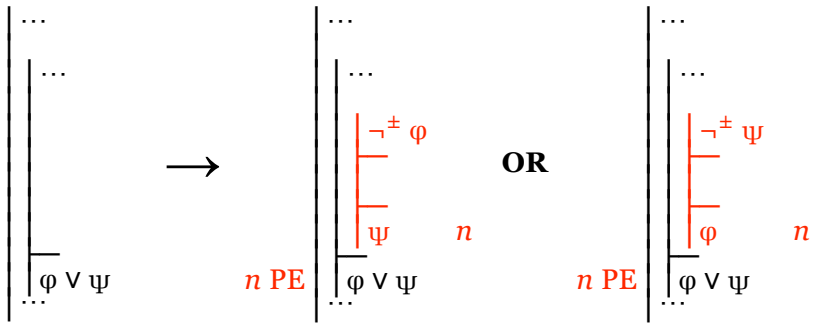


Rules from chapter 4

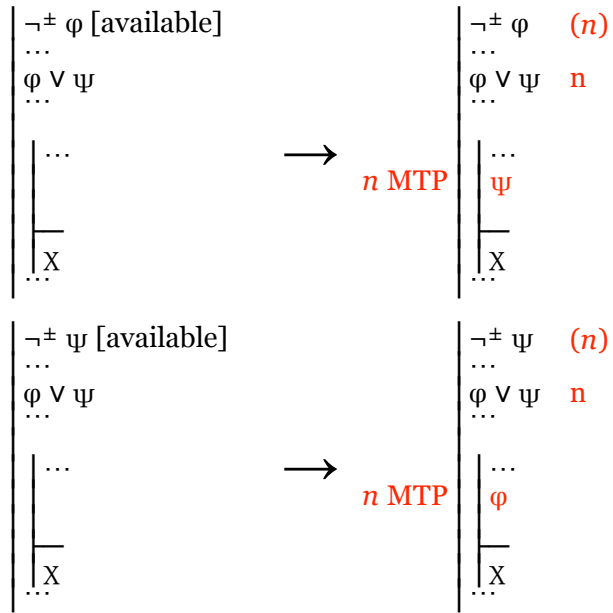
Proof by Cases (PC)



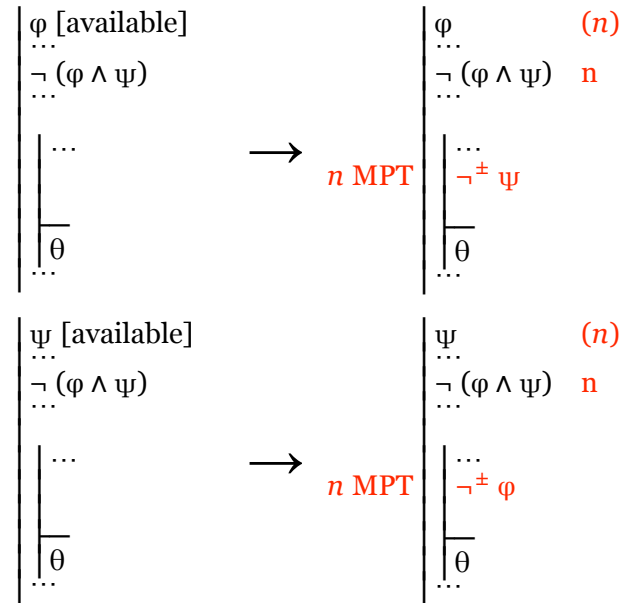
Proof of Exhaustion (PE)



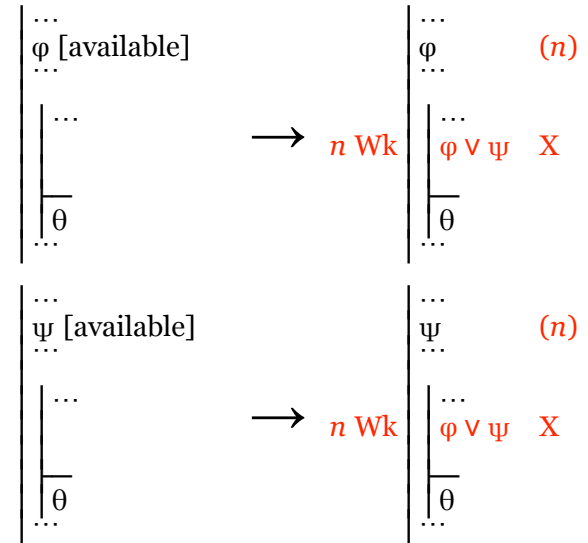
Modus Tollendo Ponens (MTP)



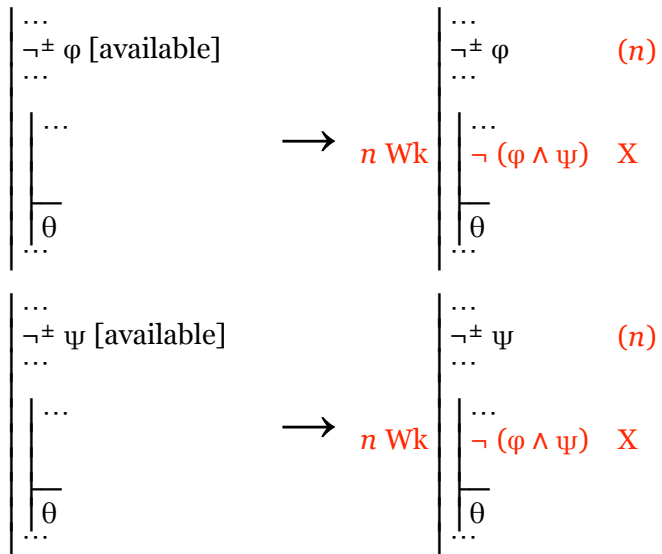
Modus Ponendo Tollens (MPT)



Weakening (Wk)

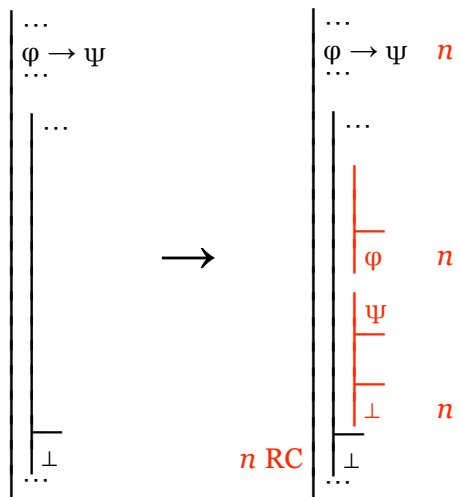


Weakening (Wk)

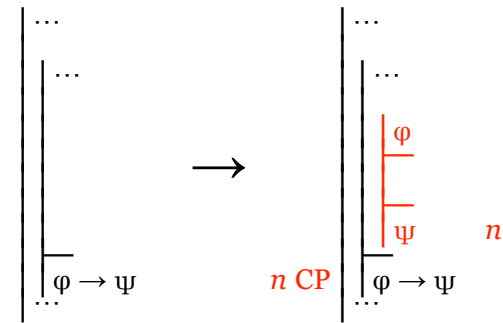


Rules from chapter 5

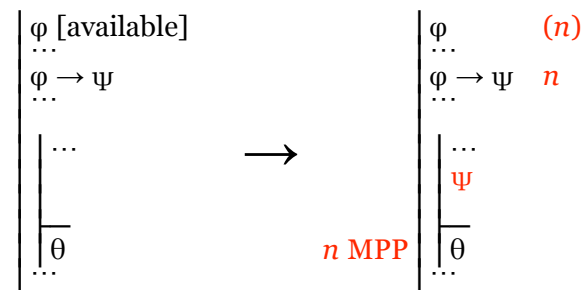
Rejecting a Conditional (RC)



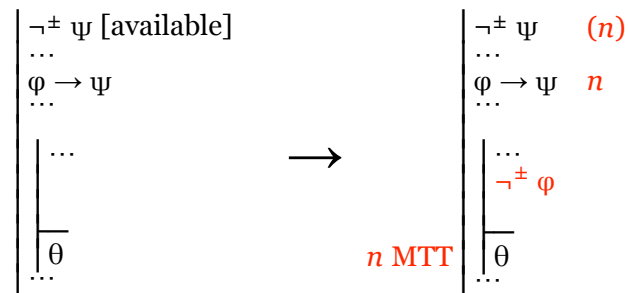
Conditional Proof (CP)



Modus Ponendo Ponens (MPP)



Modus Tollendo Tollens (MTT)



Weakening (Wk)

$$\left| \begin{array}{c} \psi \text{ [available]} \\ \dots \\ \dots \\ \hline \theta \\ \dots \end{array} \right| \rightarrow n \text{ Wk} \left| \begin{array}{c} \psi \\ \dots \\ \dots \\ \hline \theta \\ \dots \end{array} \right| \begin{array}{l} (n) \\ \varphi \rightarrow \psi \text{ X} \end{array}$$

Weakening (Wk)

$$\left| \begin{array}{c} \neg^\pm \varphi \text{ [available]} \\ \dots \\ \dots \\ \hline \theta \\ \dots \end{array} \right| \rightarrow n \text{ Wk} \left| \begin{array}{c} \neg^\pm \varphi \\ \dots \\ \dots \\ \hline \theta \\ \dots \end{array} \right| \begin{array}{l} (n) \\ \varphi \rightarrow \psi \text{ X} \end{array}$$

Rules from chapter 6

Equated Co-aliases (EC)

$$\left| \begin{array}{c} \dots \\ [\tau \text{ and } u \text{ are co-aliases}] \\ \dots \\ \dots \\ \hline \tau = u \\ \dots \end{array} \right| \rightarrow n \text{ EC} \left| \begin{array}{c} \dots \\ [\tau \text{ and } u \text{ are co-aliases}] \\ \dots \\ \dots \\ \hline \tau = u \\ \dots \end{array} \right|$$

Distinguished Co-aliases (DC)

$$\left| \begin{array}{c} \dots \\ [\tau \text{ and } u \text{ are co-aliases}] \\ \dots \\ \neg \tau = u \\ \dots \\ \dots \\ \hline \perp \\ \dots \end{array} \right| \rightarrow n \text{ DC} \left| \begin{array}{c} \dots \\ [\tau \text{ and } u \text{ are co-aliases}] \\ \dots \\ \neg \tau = u \\ \dots \\ \dots \\ \hline \perp \\ \dots \end{array} \right| \begin{array}{l} (n) \\ \text{red dot} \end{array}$$

QED given equations (QED=)

$$\left| \begin{array}{c} \dots \\ [\tau_1 \dots \tau_n \text{ and } u_1 \dots u_n \\ \text{are co-alias' series}] \\ \dots \\ P\tau_1 \dots \tau_n \\ \dots \\ \dots \\ \hline P u_1 \dots u_n \\ \dots \end{array} \right| \rightarrow n \text{ QED=} \left| \begin{array}{c} \dots \\ [\tau_1 \dots \tau_n \text{ and } u_1 \dots u_n \\ \text{are co-alias' series}] \\ \dots \\ P\tau_1 \dots \tau_n \\ \dots \\ \dots \\ \hline P u_1 \dots u_n \\ \dots \end{array} \right| \begin{array}{l} (n) \\ \text{red dot} \end{array}$$

Note: Two series of terms are co-alias series when their corresponding members are co-aliases.

Non-contradiction given equations (Nc=)

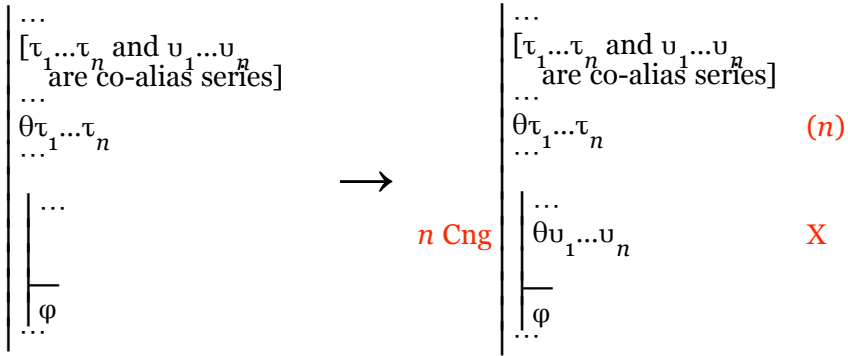
$$\left| \begin{array}{c} \dots \\ [\tau_1 \dots \tau_n \text{ and } u_1 \dots u_n \\ \text{are co-alias' series}] \\ \dots \\ \neg P\tau_1 \dots \tau_n \\ \dots \\ P u_1 \dots u_n \\ \dots \\ \dots \\ \hline \perp \\ \dots \end{array} \right| \rightarrow n \text{ Nc=} \left| \begin{array}{c} \dots \\ [\tau_1 \dots \tau_n \text{ and } u_1 \dots u_n \\ \text{are co-alias' series}] \\ \dots \\ \neg P\tau_1 \dots \tau_n \\ \dots \\ P u_1 \dots u_n \\ \dots \\ \dots \\ \hline \perp \\ \dots \end{array} \right| \begin{array}{l} (n) \\ \text{red dot} \end{array}$$

Note: Two series of terms are co-alias series when their corresponding members are co-aliases.

Co-alias Equation (CE)

$$\left| \begin{array}{c} \dots \\ [\tau \text{ and } u \text{ are co-aliases}] \\ \dots \\ \dots \\ \hline \varphi \\ \dots \end{array} \right| \rightarrow n \text{ CE} \left| \begin{array}{c} \dots \\ [\tau \text{ and } u \text{ are co-aliases}] \\ \dots \\ \dots \\ \hline \tau = u \\ \dots \\ \hline \varphi \\ \dots \end{array} \right| \text{ X}$$

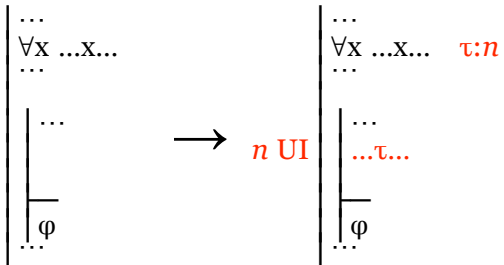
Congruence (Cng)



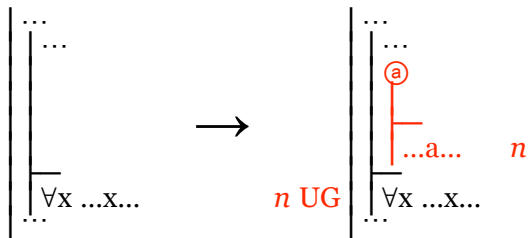
Note: θ can be an abstract, so $\theta\tau_1\dots\tau_n$ and $\theta u_1\dots u_n$ are any formulas that differ only in the occurrence of terms and in which the corresponding terms are co-aliases.

Rules from chapter 7

Universal Instantiation (UI)

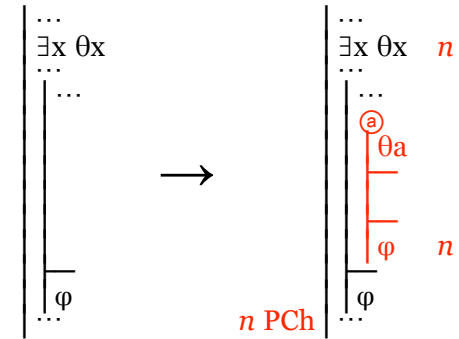


Universal Generalization (UG)

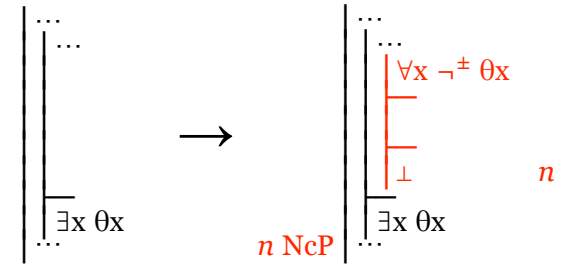


Rules from chapter 8

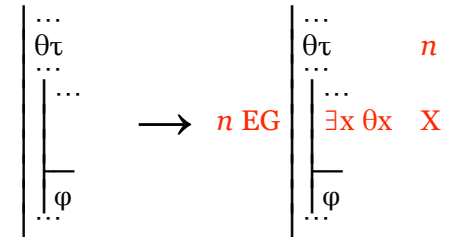
Proof by Choice (PCh)



Non-constructive Proof (NcP)



Existential Generalization (EG)



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