

5.4.s. Summary

The **law for the conditional as a premise** applies only to *reductio* arguments and provides a way of **rejecting** a conditional by deriving its antecedent ϕ from the premises and reducing its consequent to absurdity given the premises. The corresponding derivation rule is **Rejecting a Conditional (RC)**.

This rule reflects the fact that a conditional is false when its antecedent is true and its consequent is false. The rules of **Weakening (Wk)** that have conditionals as conclusions reflect the fact that a conditional is true if its consequent is and also if its antecedent is false.

With these rules, the system of derivations for truth-functional logic is complete. It consists of the fundamental rules for developing gaps by exploiting resources or planning for goals, two rules each for negations, conjunctions, disjunctions, and conditionals along with a rule to plan for atomic sentences. There are the same four rules for closing gaps we had as of 3.2, and we now also have a set of four detachment rules that provide alternative ways of exploiting weak truth-functional compounds. These rules form the basic system; and all are progressive. In addition, there is a group of rules that are not necessarily progressive although they are sound and safe—the attachment rules and the general rule LFR for introducing lemmas in *reductio* arguments.

Rules for developing gaps			Rules for closing gaps	
	for resources	for goals	when to close	rule
atomic sentence		IP	the goal is also a resource	QED
negation $\neg \phi$ (if ϕ is not atomic and the goal is \perp)	CR	RAA	sentences ϕ and $\neg \phi$ are resources & the goal is \perp	Nc
conjunction $\phi \wedge \psi$	Ext	Cnj	\top is the goal	ENV
disjunction $\phi \vee \psi$	PC	PE	\perp is a resource	EFQ
conditional $\phi \rightarrow \psi$ (if the goal is \perp)	RC	CP		
<i>Basic system</i>				
<i>Detachment rules (optional)</i>				
main resource	auxiliary resource	rule	<i>Attachment rules</i>	
	ϕ	MPP	added resource	rule
$\phi \rightarrow \psi$	$\bar{\psi}$	MTT	$\phi \wedge \psi$	Adj
$\phi \vee \psi$	$\bar{\phi}$ or $\bar{\psi}$	MTP	$\phi \rightarrow \psi$	Wk
$\neg (\phi \wedge \psi)$	ϕ or ψ	MPT	$\phi \vee \psi$	Wk
			$\neg (\phi \wedge \psi)$	Wk
			<i>Rule for lemmas</i>	
			prerequisite	rule
			the goal is \perp	LFR
				<i>Added rules (optional)</i>

As in the earlier tables of this form, the names of the rules are links to places where they are actually stated.

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